

1 A Numerical Analyst's Tubular Neighborhood Theorem

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4 **Abstract**

5 For complex analytic manifolds in Euclidean spaces, this paper establishes
6 a version of Tubular Neighborhood Theorem that is specifically formulated for
7 the applications in regularization of ill-posed algebraic problems for numer-
8 ical computation. The theorem and the geometric analysis are indispensable
9 for achieving accurate and stable numerical solutions of those highly sensitive
10 problems using floating point arithmetic even if the data may be perturbed.

11 **1 Introduction**

12 When a scientific computing problem is attempted, there are at least three problems
13 lurking around: The underlying problem that is intended but may be hidden, the
14 given problem embodied in the available data, and the actual problem that ends up
15 being solved. Those problems are not exactly the same but the differences are usually
16 too small to be noticed until certain unexpected solutions emerge.

17 Think about a seemingly simple root-finding problem for a given polynomial

$$\tilde{p}(x) = x^5 + 4.667 x^4 + 8.667 x^3 + 8 x^2 + 3.667 x + .6667, \tag{1}$$

18 What is really intended may be roots of the hidden polynomial

$$p(x) = x^5 + \frac{14}{3} x^4 + \frac{26}{3} x^3 + 8 x^2 + \frac{11}{3} x + \frac{2}{3} \equiv (x + 1)^4 (x + \frac{2}{3}) \tag{2}$$

19 for which $\tilde{p}(x)$ serves as empirical data. Sending the data polynomial $\tilde{p}(x)$ to any
20 commercial software such as Matlab for its roots, the results such as

-1.1211497 + 0.1032009i, -1.1211497 - 0.1032009i, -0.8863741 + 0.1451301i, -0.8863741 - 0.1451301i, -0.6519516

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21 are far from expected even though they match exact roots of \tilde{p} for the first three
 22 to four digits. The actual problem being solved is the exact root collection of the
 23 polynomial

$$\begin{aligned}\hat{p}(x) &= (x_{+1.1211497 - 0.1032009i})(x_{+1.1211497 + 0.1032009i}) \cdots (x_{+0.6519516}) \\ &\approx x^5 + 4.666999 x^4 + 8.666997 x^3 + 7.999998 x^2 + 3.6669996 x + .6667001\end{aligned}$$

24 All three problems are nearby as we expected, but the computing results bear no
 25 resemblance to the intended solutions $-1, -\frac{2}{3}$, as shown in Figure 1.

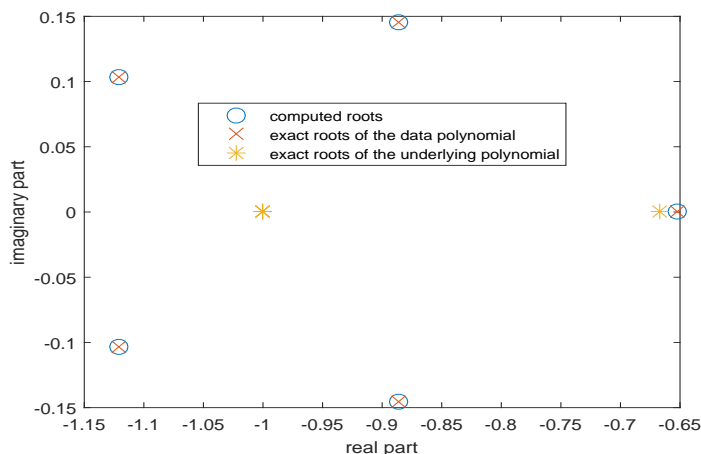


Figure 1: Computed roots and roots of the data polynomial \tilde{p} are far from the roots of the underlying polynomial p due to hypersensitivity.

26 Dubbed as one of the “Maxims about numerical mathematics, science, computers
 27 and life on earth” [13]: “If the answer is highly sensitive to perturbations, you have
 28 probably asked the wrong question”. Attempt in solving the root-finding problem
 29 of $\tilde{p}(x)$ in conventional sense is apparently such a wrong problem since a tiny
 30 perturbation from $p(x)$ to $\hat{p}(x)$ alters the solution seemingly completely.

31 There is also a fourth problem nearby. Wouldn’t it make more sense to seek what
 32 we shall establish as the *regularized roots* of $\tilde{p}(x)$ and its *regularized factorization*

$$1.000068 (x + 0.99994)^4 (x + 0.66682),$$

33 particularly after proving such roots and factorization are existent, unique, Lipschitz
 34 continuous with respect to the data \tilde{p} , approximate those of the hidden polynomial
 35 $p(x)$, and can be computed accurately using floating point arithmetic from empirical
 36 data [15, 16]?

37 Those problems with extremely high sensitivities beyond finite bounds are known
 38 to be *ill-posed problems*, a concept attributed to Hadamard. As it turns out, such
 39 problems are abundant in scientific computing. Besides root-finding, even some of
 40 the most basic algebraic problems are ill-posed, such as matrix ranks and subspaces,

41 polynomial greatest common divisors and factorizations, defective eigenvalues and
42 Jordan Canonical Forms, etc. Those are the problems we will encounter inevitably
43 in scientific computing.

44 Are the roots of aforementioned $p(x)$ really sensitive to data perturbations as alleged?

45 In a legendary technical report began circulating in 1972 [9] but never published,
46 Kahan argues that it is a “misconception” to consider multiple roots hypersensitive
47 to data perturbations, and he proves that the sensitivity of a root is finitely bounded
48 if the data perturbation is constrained such that the multiplicity is preserved and thus
49 the identity of the root remains intact. Kahan further argues that, albeit without
50 proof, the ill-posed problems of a certain structure form what he called a “pejorative
51 manifold”. A problem is sensitive only if the data point is *near* the manifold but not
52 necessarily so *on* it.

53 Hypersensitive problems frequently form complex analytic manifolds, such as those
54 problems mentioned above. Furthermore, those manifolds entangle to form strata
55 in which every manifold is embedded in the closure of some other manifolds of lower
56 codimensions. As a result, the hypersensitivity of the problem is not random but
57 directional: The singularity of a problem can only be decreased and never increases by
58 a tiny perturbation. Using the root-finding example above, the roots of $p(x)$ changes
59 in one and only one direction: The multiplicity 4 reduces by tiny perturbations.

60 The phenomenon that certain hypersensitive problems form a manifold, rather than
61 merely an algebraic variety, is crucial in the analysis, regularization and numerical
62 solution of those problems thanks to the existence of the *tubular neighborhood* that is
63 considered “one of the most useful notions in the theory of differentiable manifolds”
64 [5]. When the problem data are given as empirical and subject to small perturbations
65 due to measurement and round-off, we are having a point near the complex analytic
66 manifold in the problem space. Assuming the data are reasonably accurate so that the
67 point remains in the tubular neighborhood of the manifold, the Tubular Neighborhood
68 Theorem ensures the projection from the empirical data point to the manifold to be
69 existent, unique and Lipschitz continuous, making the projection a well-posed least
70 squares problem that is accurately solvable in numerical computation.

71 Geometric theories and insights have been applied in numerical analysis effectively
72 in the works such as Sommese, Verschelde and Wampler[11], Corless, Galligo, Kot-
73 sireas and Watt [3], Edelman, Elmroth and Kågström [7, 8], Lippert and Edelman
74 [10], Demmel and Edelman [4], Absil, Mahony and Sepulchre [1], Edelman, Arias and
75 Smith [6], and so on. However, the tremendous advantage of the tubular neighbor-
76 hoods has apparently not yet been utilized in numerical analysis. Specifically aimed
77 at the application of solving hypersensitive algebraic problems in this paper, we es-
78 tablish a version of Tubular Neighborhood Theorem for complex analytic manifold in
79 Euclidean spaces isometrically isomorphic to \mathbb{C}^n , and provide an elementary proof
80 with a dose of flavor in numerical analysis. The theorem and the proof are not in-
81 tended to impress differential topologists but to fill a gap in the regularization theory

82 of hypersensitive problems and to complete the works of numerical factorization and
83 numerical greatest common divisors of polynomials.

84 This paper attempts to provide a synergy between numerical analysis and differential
85 topology. Besides establishing the version of Tubular Neighborhood Theorem for
86 numerical analysis applications, we outline the geometric analysis of ill-posed prob-
87 lems in a series of observations, the regularization theory based on the geometry and
88 the resulting strategy for solving the regularized problem accurately in numerical
89 computation even if the data may be perturbed. These geometric analyses set the
90 foundation for our software package NACLAB¹ for numerical algebraic computing.

91 2 Preliminaries

92 The space of n -dimensional vectors of complex numbers is denoted by \mathbb{C}^n as
93 a *topological vector space*. As a vector space, it is closed to addition and scalar
94 multiplication. In the sense of a topological space, its topology is derived from the
95 Euclidean norm $\|\cdot\|_2$.

96 We consider *finite-dimensional normed vector spaces* denoted by, say \mathcal{V} , \mathcal{W} and so
97 on. The norm $\|\mathbf{v}\|$ is understood from the context as *the* norm in the space where
98 \mathbf{v} belongs. Such a vector space \mathcal{V} is isomorphic to \mathbb{C}^n where $n = \dim(\mathcal{V})$, the
99 dimension of \mathcal{V} .

100 A mapping $F : \Omega \subset \mathbb{C}^n \rightarrow \mathbb{C}^m$ is *holomorphic* if every one of the m components
101 of F is a holomorphic function in every one of the n variables in the open subset
102 Ω of \mathbb{C}^n . We can designate a variable name, say \mathbf{z} , for F and denote F as
103 $\mathbf{z} \mapsto F(\mathbf{z})$. In that case the Jacobian matrix of F at any particular $\mathbf{z}_0 \in \Omega$ is
104 denoted by $F_{\mathbf{z}}(\mathbf{z}_0)$.

105 Let \mathcal{V} and \mathcal{W} be finite-dimensional normed vector spaces isomorphic to \mathbb{C}^n and
106 \mathbb{C}^m respectively via isomorphisms $\psi_{\mathcal{V}} : \mathcal{V} \rightarrow \mathbb{C}^n$ and $\psi_{\mathcal{W}} : \mathcal{W} \rightarrow \mathbb{C}^m$. Assume \mathbf{g}
107 is a mapping from an open subset Σ of \mathcal{V} to \mathcal{W} with a representation $\mathbf{z} \mapsto G(\mathbf{z})$
108 where $G : \psi_{\mathcal{V}}(\Sigma) \subset \mathbb{C}^n \rightarrow \mathbb{C}^m$ such that $\mathbf{g} = \psi_{\mathcal{W}}^{-1} \circ G \circ \psi_{\mathcal{V}}$. We say \mathbf{g} is
109 holomorphic in Σ if G is holomorphic in $\psi_{\mathcal{V}}(\Sigma)$. Denoting the variable of \mathbf{g}
110 as, say \mathbf{v} , the *Jacobian* of \mathbf{g} at any particular $\mathbf{v}_0 \in \Sigma$ is defined as the linear
111 transformation $\mathbf{g}_{\mathbf{v}}(\mathbf{v}_0) : \mathcal{V} \rightarrow \mathcal{W}$ in the form of

$$\mathbf{v} \mapsto \mathbf{g}_{\mathbf{v}}(\mathbf{v}_0)(\mathbf{v}) = \psi_{\mathcal{W}}^{-1} \circ G_{\mathbf{z}}(\mathbf{z}_0) \circ \psi_{\mathcal{V}}(\mathbf{v})$$

112 where $\mathbf{z}_0 = \psi_{\mathcal{V}}(\mathbf{v}_0)$. The Jacobian $\mathbf{g}_{\mathbf{v}}(\mathbf{v}_0)$ so defined as a linear transformation is
113 invariant under change of bases and the respective isomorphisms. Let $G_{\mathbf{z}}(\mathbf{z}_0)^{\mathbf{H}}$ and
114 $G_{\mathbf{z}}(\mathbf{z}_0)^{\dagger}$ be the Hermitian transpose and the Moore-Panrose inverse of the Jacobian
115 matrix $G_{\mathbf{z}}(\mathbf{z}_0)$ respectively. If we further assume the isomorphisms $\psi_{\mathcal{V}}$ and $\psi_{\mathcal{W}}$

¹<http://homepages.neiu.edu/~nac1ab>

116 are isometric so that $\|\psi_{\mathcal{V}}(\mathbf{v})\|_2 = \|\mathbf{v}\|$ and $\|\psi_{\mathcal{W}}(\mathbf{w})\|_2 = \|\mathbf{w}\|$ for all $\mathbf{v} \in \mathcal{V}$ and
 117 $\mathbf{w} \in \mathcal{W}$, then the Hermitian transpose and the Moore-Panrose inverse of $\mathbf{g}_{\mathbf{v}}(\mathbf{v}_0)$ are
 118 well-defined as linear transformations $\mathbf{g}_{\mathbf{v}}(\mathbf{v}_0)^{\text{H}}$ and $\mathbf{g}_{\mathbf{v}}(\mathbf{v}_0)^{\dagger}$ respectively from \mathcal{W}
 119 to \mathcal{V} in the forms of

$$\begin{aligned} \mathbf{g}_{\mathbf{v}}(\mathbf{v}_0)^{\text{H}} &= \psi_{\mathcal{V}}^{-1} \circ G_{\mathbf{z}}(\mathbf{z}_0)^{\text{H}} \circ \psi_{\mathcal{W}} \\ \mathbf{g}_{\mathbf{v}}(\mathbf{v}_0)^{\dagger} &= \psi_{\mathcal{V}}^{-1} \circ G_{\mathbf{z}}(\mathbf{z}_0)^{\dagger} \circ \psi_{\mathcal{W}} \end{aligned}$$

120 that are invariant under any change of bases resulting in isometric isomorphisms.
 121 The holomorphic mapping $\mathbf{v} \mapsto \mathbf{g}(\mathbf{v})$ is said to be *non-degenerate* at \mathbf{v}_0 if the
 122 Jacobian $\mathbf{g}_{\mathbf{v}}(\mathbf{v}_0)$ is an injective linear transformation.

123 The notion of holomorphic mapping extends beyond open domains. A mapping
 124 $\mathbf{f} : \Pi \subset \mathcal{V} \rightarrow \mathcal{W}$ is holomorphic in a non-open domain Π if there is an open subset
 125 Ω of \mathcal{V} containing Π and a holomorphic mapping \mathbf{g} defined in Ω such that
 126 $\mathbf{f}(\mathbf{z}) \equiv \mathbf{g}(\mathbf{z})$ for all $\mathbf{z} \in \Pi$.

127 3 Complex analytic manifolds

128 In this paper, applications will be studied and computations will be conducted in
 129 finite-dimensional normed vector spaces in which we shall consider complex analytic
 130 manifolds in topological sense.

131 A finite-dimensional normed vector space \mathcal{V} has an obvious topology in which an
 132 ε -neighborhood $N_{\varepsilon}(\mathbf{v}) := \{\mathbf{z} \in \mathcal{V} \mid \|\mathbf{z} - \mathbf{v}\| < \varepsilon\}$ of any $\mathbf{v} \in \mathcal{V}$ and $\varepsilon > 0$ is
 133 an open set along with unions of any collection of those neighborhoods. The vector
 134 space \mathcal{V} equipped with this topology becomes a topological vector space. A subset
 135 $\Sigma \subset \mathcal{V}$ is a (topological) subspace of \mathcal{V} with the subspace topology in which Λ is
 136 an open set if and only if there is an open set Ω in \mathcal{V} such that $\Lambda = \Omega \cap \Sigma$.

137 The standard definition of manifolds in differential topology is more abstract and
 138 serves broader purposes than what we need for our applications in numerical solutions
 139 of ill-posed algebraic problems. We restrict our elaborations to the type of complex
 140 analytic manifolds in Euclidean spaces. They are, in fact, submanifolds of finite-
 141 dimensional normed vector spaces.

142 **Definition 3.1 (Complex Analytic Manifold)** *A topological subspace Π of a*
 143 *finite-dimensional normed vector space \mathcal{V} over \mathbb{C} is called a complex analytic*
 144 *manifold of dimension m if, for every point $\mathbf{x} \in \Pi$, there is an open neighborhood*
 145 *Σ of \mathbf{x} in \mathcal{V} and a holomorphic mapping ϕ from $\Sigma \cap \Pi$ into an open subset*
 146 *Λ of \mathbb{C}^m with a holomorphic inverse ψ such that $\phi \circ \psi(\mathbf{z}) \equiv \mathbf{z}$ for all $\mathbf{z} \in \Lambda$*
 147 *and $\psi \circ \phi(\mathbf{y}) \equiv \mathbf{y}$ for all $\mathbf{y} \in \Sigma \cap \Pi$.*

148 The dimension deficit $\dim(\mathcal{V}) - m$ is called the *codimension* of Π in \mathcal{V} denoted
 149 by $\text{codim}(\Pi)$. The vector space \mathcal{V} itself is a complex analytic manifold in \mathcal{V} in
 150 which every complex analytic manifold is a submanifold.

151 **Example 1 (Rank Manifold)** [4] In the vector space $\mathbb{C}^{m \times n}$ of $m \times n$ matrices
 152 of complex entries with the Frobenius norm, the subset

$$\mathcal{C}_r^{m \times n} = \{A \in \mathbb{C}^{m \times n} \mid \text{rank}(A) = r\} \quad (3)$$

153 is a complex analytic manifold of codimension $(m - r)(n - r)$. Here $\text{rank}(A)$
 154 denotes the rank of the matrix A . The proof is quite straightforward: For any
 155 matrix $A \in \mathcal{C}_r^{m \times n}$, there are r linearly independent columns forming a submatrix
 156 $X \in \mathbb{C}^{m \times r}$. The submatrix of the remaining $n - r$ columns can be written as XY
 157 where $Y \in \mathbb{C}^{r \times (n-r)}$. The mapping $A \mapsto (X, Y)$ is holomorphic and invertible
 158 from a neighborhood of A in Π into the vector space $\mathbb{C}^{m \times r} \times \mathbb{C}^{r \times (n-r)}$ that is
 159 isometrically isomorphic to $\mathbb{C}^{mn - (m-r)(n-r)}$.

160 **Example 2 (Univariate Factorization Manifold)** [14] Let the vector space

$$\mathbb{P}_n = \{a_0 + a_1 x + \cdots + a_n x^n \mid a_0, \dots, a_n \in \mathbb{C}\}$$

161 with the norm $\|a_0 + a_1 x + \cdots + a_n x^n\| = \sqrt{|a_0|^2 + \cdots + |a_n|^2}$, making \mathbb{P}_n
 162 isometrically isomorphic to \mathbb{C}^{n+1} . For any positive integers $\ell_1 + \cdots + \ell_k = n$, the
 163 subset

$$\mathcal{F}_{\ell_1 \dots \ell_k} = \{\alpha (x - z_1)^{\ell_1} \cdots (x - z_k)^{\ell_k} \in \mathbb{P}_n \mid \alpha, z_1, \dots, z_k \in \mathbb{C}, z_i \neq z_j \text{ for } i \neq j\} \quad (4)$$

164 is a complex analytic manifold of codimension $n - k$ in \mathbb{P}_n .

165 **Example 3 (Univariate GCD Manifold)** [18] The product space $\mathbb{P}_m \times \mathbb{P}_n$ with
 166 the norm $\|(p, q)\| = \sqrt{\|p\|^2 + \|q\|^2}$ is isometrically isomorphic to \mathbb{C}^{m+n+2} . For
 167 every pair $(p, q) \in \mathbb{P}_m \times \mathbb{P}_n$, let $\text{gcd}(p, q)$ denote the greatest common divisor of p
 168 and q . The subset

$$\mathcal{P}_{m,n}^k = \{(p, q) \in \mathbb{P}_m \times \mathbb{P}_n \mid \text{deg}(p) = m, \text{deg}(q) = n, \text{deg}(\text{gcd}(p, q)) = k\}$$

169 is a complex analytic manifold of codimension k in $\mathbb{P}_m \times \mathbb{P}_n$ where $\text{deg}(\cdot)$ is the
 170 degree of any polynomial (\cdot) .

171 Not all algebraic varieties are manifolds, as shown in the following example.

172 **Example 4 (Example of a non-manifold)** Consider the subset

$$\Pi_5^2 = \{(x - \alpha)^2 q(x) \in \mathbb{P}_5 \mid q \in \mathbb{P}_3, q(\alpha) \neq 0\}$$

173 of \mathbb{P}_5 having a double root and define the mapping

$$\mathbf{f} : (\alpha, a_0, a_1, a_2, a_3) \mapsto (x - \alpha)^2 (a_0 + a_1 x + a_2 x^2 + a_3 x^3)$$

174 from \mathbb{C}^5 to \mathbb{P}_5 . The polynomial

$$\begin{aligned} p(x) &= (x - 1)^2 (x - 2)^2 (x - 3) \\ &\equiv (x - 1)^2 (-12 + 16x - 7x^2 + x^3) \\ &\equiv (x - 2)^2 (-3 + 7x - 5x^2 + x^3) \end{aligned}$$

175 in Π_5^2 is the image of both $\mathbf{f}(1, -12, 16, -7, 1)$ and $\mathbf{f}(2, -3, 7, -5, 1)$. Namely,
 176 the polynomial p is in the intersection $\mathbf{f}(\Omega_1) \cap \mathbf{f}(\Omega_2)$ of two branches of Π_5^2 where
 177 $\Omega_1, \Omega_2 \subset \mathbb{C}^5$ are small neighborhoods of $(1, -12, 16, -7, 1)$ and $(2, -3, 7, -5, 1)$
 178 respectively. As a result, the subset Π_5^2 is not a complex analytic manifold.

179 4 The Tubular Neighborhood Theorem

180 The Tubular Neighborhood Theorem can be illustrated in Figure 2: A complex ana-
 181 lytic manifold is contained in an open *tubular neighborhood* in which every point can
 182 be uniquely projected onto the manifold following a normal line. For our application
 183 of numerical solution of ill-posed problems, it is sufficient for the projection from the
 184 tubular neighborhood to the manifold to be Lipschitz continuous. Being smooth or
 185 analytic is a unnecessary luxury. The Lipschitz constant can also serve as a condition
 186 number.

187 The Tubular Neighborhood Theorem is one of the fundamental results in differential
 188 topology. Standard versions of the Tubular Neighborhood Theorem for real smooth
 189 manifolds can be found in textbooks of differential geometry. Those versions are
 190 presented in abstract forms for general purposes and do not appear directly applicable
 191 to our applications involving complex analytic manifolds. The following Theorem 4.1
 192 is a down-to-earth version of the Tubular Neighborhood Theorem for submanifolds of
 193 normed vector spaces isometrically isomorphic to \mathbb{C}^n 's and it is specifically tailored
 194 for our applications. Since we can't find such a version in the literature, we provide
 195 a numerical analysis proof using the Gauss-Newton iteration that also serves as the
 196 mechanism for calculating the projection to the manifold.

197 **Theorem 4.1 (Tubular Neighborhood Theorem)** *Let Π be a complex analytic*
 198 *manifold in a normed vector space \mathcal{V} that is isometrically isomorphic to \mathbb{C}^n . There*

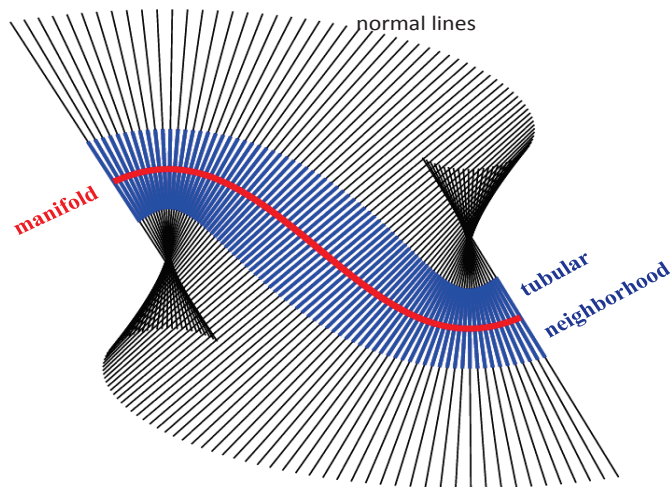


Figure 2: Illustration of the tubular neighborhood

199 *is a tubular neighborhood, namely an open subset Ω of \mathcal{V} such that $\Pi \subset \Omega$, and*
 200 *every $\mathbf{b} \in \Omega$ has a unique projection $\mathbf{x}_{\mathbf{b}} \in \Pi$ of minimum distance to Π , that is*

$$\|\mathbf{x}_{\mathbf{b}} - \mathbf{b}\| = \inf_{\mathbf{x} \in \Pi} \|\mathbf{x} - \mathbf{b}\|. \quad (5)$$

201 *Furthermore, the projection $\mathbf{b} \mapsto \mathbf{x}_{\mathbf{b}}$ from Ω to Π is locally Lipschitz continuous.*

202 We shall prove this particular version of Tubular Neighborhood Theorem in §6.

203 From the perspective of a numerical analyst, the projection from any point \mathbf{b} in
 204 the tubular neighborhood Ω to the manifold Π is the (nonlinear) least squares
 205 problem of solving the (nonlinear) system $\mathbf{f}(\mathbf{x}) = \mathbf{b}$ where $\mathbf{f} : \mathbb{C}^m \rightarrow \mathcal{V}$ is
 206 the local diffeomorphism for the manifold Π . The Tubular Neighborhood Theorem
 207 ensures the solution of such a least squares problem exists, is unique and is Lipschitz
 208 continuous.

209 5 The nonlinear least squares problem

210 Let \mathcal{V} and \mathcal{W} be normed vector spaces over \mathbb{C} that are isometrically isomorphic to
 211 \mathbb{C}^n and \mathbb{C}^m respectively with $m > n$. For any holomorphic mapping $\mathbf{x} \mapsto \mathbf{f}(\mathbf{x})$
 212 from an open subset Ω of \mathcal{V} to \mathcal{W} , consider the equation

$$\mathbf{f}(\mathbf{x}) = \mathbf{b}. \quad (6)$$

213 Since the image $\mathbf{f}(\Omega)$ is some kind of surface of dimension at most n in the m -
 214 dimensional space \mathcal{W} with $m > n$, conventional solutions do not exist for the
 215 equation (6) except in special cases where $\mathbf{b} \in \mathbf{f}(\Omega)$. Instead, we seek a *least*
 216 *squares solution* $\mathbf{x}_* \in \Omega$ such that

$$\|\mathbf{f}(\mathbf{x}_*) - \mathbf{b}\|^2 = \min_{\mathbf{x} \in \Omega} \|\mathbf{f}(\mathbf{x}) - \mathbf{b}\|^2.$$

217 In other words, we seek \mathbf{x}_* so that $\mathbf{f}(\mathbf{x}_*)$ is the projection of \mathbf{b} on to the surface
 218 $\mathbf{f}(\Omega)$ so that the minimal distance from \mathbf{b} to the surface is attained.

219 It is a straightforward verification that such a least squares solution is a *critical point*
 220 for the equation (6), namely

$$\mathbf{f}_x(\mathbf{x}_*)^H (\mathbf{f}(\mathbf{x}_*) - \mathbf{b}) = \mathbf{0}. \quad (7)$$

221 The interpretation of (7) is quite clear: At the least squares solution \mathbf{x}_* , the vector
 222 $\mathbf{f}(\mathbf{x}_*) - \mathbf{b}$ must be orthogonal to the tangent plane of the surface $\mathbf{f}(\Omega)$ at the point
 223 $\mathbf{f}(\mathbf{x}_*)$.

224 An effective method for finding the least squares solution of (6) is the *Gauss-Newton*
 225 *iteration*

$$\mathbf{x}_{k+1} = \mathbf{x}_k - \mathbf{f}_x(\mathbf{x}_k)^\dagger (\mathbf{f}(\mathbf{x}_k) - \mathbf{b}) \quad \text{for } k = 0, 1, \dots \quad (8)$$

226 The Gauss-Newton iteration (8) is locally convergent requiring the initial iterate \mathbf{x}_0
 227 to be sufficiently near the least squares solution \mathbf{x}_* and the residual $\|\mathbf{f}(\mathbf{x}_*) - \mathbf{b}\|$ to
 228 be sufficiently small. The following lemma provides detailed convergence conditions
 229 in Kantorovich style.

230 **Lemma 5.1** [17] *Let \mathcal{V} and \mathcal{W} be finite-dimensional normed vector spaces iso-*
 231 *metrically isomorphic to \mathbb{C}^n and \mathbb{C}^m respectively. Assume $\mathbf{x} \mapsto \mathbf{f}(\mathbf{x})$ is a*
 232 *holomorphic mapping from an open subset $\Omega \subset \mathcal{V}$ to \mathcal{W} with a critical point*
 233 *$\mathbf{x}_* \in \Omega$ of the system $\mathbf{f}(\mathbf{x}) = \mathbf{b}$. Assume $\Lambda \subset \Omega$ is an open neighborhood of \mathbf{x}_**
 234 *so that \mathbf{f} is non-degenerate in Λ and there exist constants $\zeta, \gamma > 0$ such that*

$$\|\mathbf{f}_x(\mathbf{z})^\dagger\| \leq \zeta, \quad (9)$$

$$\|\mathbf{f}(\mathbf{z}) - \mathbf{f}(\tilde{\mathbf{z}}) - \mathbf{f}_x(\tilde{\mathbf{z}})(\mathbf{z} - \tilde{\mathbf{z}})\| \leq \gamma \|\mathbf{z} - \tilde{\mathbf{z}}\|^2 \quad (10)$$

235 for all $\mathbf{z}, \tilde{\mathbf{z}} \in \Lambda$. Further assume $\|\mathbf{f}(\mathbf{x}_*) - \mathbf{b}\|$ is small so that

$$\|(\mathbf{f}_x(\mathbf{z})^\dagger - \mathbf{f}_x(\mathbf{x}_*)^\dagger)(\mathbf{f}(\mathbf{x}_*) - \mathbf{b})\| \leq \sigma \|\mathbf{z} - \mathbf{x}_*\| \quad (11)$$

236 for a constant $\sigma < 1$ at every $\mathbf{z} \in \Lambda$. Then for all $\mathbf{x}_0 \in \Lambda$ satisfying

$$\|\mathbf{x}_0 - \mathbf{x}_*\| < \frac{1-\sigma}{\zeta\gamma} \quad (12)$$

237 and $\{\mathbf{x} \in \mathcal{V} \mid \|\mathbf{x} - \mathbf{x}_*\| < \|\mathbf{x}_0 - \mathbf{x}_*\|\} \subset \Lambda$, the Gauss-Newton iteration (8) is well
 238 defined in Λ , converges to \mathbf{x}_* , and satisfies

$$\|\mathbf{x}_{k+1} - \mathbf{x}_*\| \leq (\sigma + \zeta\gamma \|\mathbf{x}_k - \mathbf{x}_*\|) \|\mathbf{x}_k - \mathbf{x}_*\|$$

239 for $k = 0, 1, \dots$ with $\mu = \sigma + \zeta\gamma \|\mathbf{x}_0 - \mathbf{x}_*\| < 1$.

240 If it is convergent, the Gauss-Newton iteration converges at a linear rate in general
 241 and at a quadratic rate if the residual $\|\mathbf{f}(\mathbf{x}_*) - \mathbf{b}\| = 0$ in the cases where the least
 242 squares solution is a conventional solution.

243 6 An elementary proof of the Tubular Neighbor- 244 hood Theorem

245 **Proof of Theorem 4.1.** Let \mathbf{z}_0 be any particular point in Π . Since Π is a complex
 246 analytic manifold in \mathcal{V} , there is a bounded neighborhood \mathcal{M} of \mathbf{z}_0 in \mathcal{V} , an open
 247 subset \mathcal{N} of \mathbb{C}^m where $m = \dim(\Pi)$, and a holomorphic mapping $\mathbf{x} \mapsto \mathbf{f}(\mathbf{x})$
 248 from \mathcal{N} to $\mathcal{M} \cap \Pi$ with a holomorphic inverse \mathbf{f}^{-1} . Denote $\mathbf{f}^{-1}(\mathbf{z}_0) = \mathbf{x}_0$. We
 249 can assume \mathcal{M} is sufficiently small and there are constants $\zeta, \gamma > 0$ such that

$$\begin{aligned} \|\mathbf{f}_x(\mathbf{y})^\dagger\| &\leq \zeta \\ \|\mathbf{f}(\mathbf{y}) - \mathbf{f}(\tilde{\mathbf{y}}) - \mathbf{f}_x(\tilde{\mathbf{y}})(\mathbf{y} - \tilde{\mathbf{y}})\| &\leq \gamma \|\mathbf{y} - \tilde{\mathbf{y}}\|^2 \end{aligned}$$

250 for all $\mathbf{y}, \tilde{\mathbf{y}} \in \mathcal{N}$. Let $N_\varepsilon(\mathbf{x}_0) = \{\mathbf{x} \in \mathbb{C}^m \mid \|\mathbf{x} - \mathbf{x}_0\| < \varepsilon\}$ and assume $\varepsilon > 0$ to be
 251 sufficiently small so that $N_\varepsilon(\mathbf{x}_0) \subset \mathcal{N}$ and $\|\mathbf{y} - \tilde{\mathbf{y}}\| < \frac{1}{2\zeta\gamma}$ for any $\mathbf{y}, \tilde{\mathbf{y}} \in N_\varepsilon(\mathbf{x}_0)$.
 252 Set $\varepsilon' = \frac{1}{3}\varepsilon$.

253 Let $\tilde{\mathcal{M}}$ be an open subset of \mathcal{M} with $\mathbf{z}_0 \in \tilde{\mathcal{M}}$ and $\mathbf{f}^{-1}(\tilde{\mathcal{M}} \cap \Pi) \subset N_{\varepsilon'}(\mathbf{x}_0)$. Then
 254 there is an open subset $\hat{\mathcal{M}}$ of $\tilde{\mathcal{M}}$ with $\mathbf{z}_0 \in \hat{\mathcal{M}}$ such that, for every $\mathbf{b} \in \hat{\mathcal{M}}$,

$$\|\mathbf{b} - \mathbf{z}_0\| < \|\mathbf{b} - \mathbf{d}\| \quad \text{for every } \mathbf{d} \in \mathcal{V} \setminus \tilde{\mathcal{M}}. \quad (13)$$

255 Since $\overline{\tilde{\mathcal{M}}} \cap \Pi$ is compact, there is a $\mathbf{z}_b \in \overline{\tilde{\mathcal{M}}} \cap \Pi$ such that

$$\|\mathbf{b} - \mathbf{z}_b\| = \min_{\mathbf{u} \in \overline{\tilde{\mathcal{M}}} \cap \Pi} \|\mathbf{b} - \mathbf{u}\| \leq \|\mathbf{b} - \mathbf{z}_0\|$$

256 and thus $\mathbf{z}_b \in \tilde{\mathcal{M}} \cap \Pi$ due to (13). We have thus proved that, for every $\mathbf{b} \in \hat{\mathcal{M}}$,
 257 there exists an $\mathbf{x}_b \in N_{\varepsilon'}(\mathbf{x}_0)$ such that $\mathbf{f}(\mathbf{x}_b) = \mathbf{z}_b$ and

$$\|\mathbf{f}(\mathbf{x}_b) - \mathbf{b}\| = \min_{\mathbf{u} \in \Pi} \|\mathbf{b} - \mathbf{u}\| = \min_{\mathbf{x} \in \mathcal{N}} \|\mathbf{f}(\mathbf{x}) - \mathbf{b}\|$$

258 and such an \mathbf{x}_b can only occur in $N_{\varepsilon'}(\mathbf{x}_0)$. In other words, the least squares
 259 solution to the equation $\mathbf{f}(\mathbf{x}) = \mathbf{b}$ and $\mathbf{z}_b = \mathbf{f}(\mathbf{x}_b)$ exists for every $\mathbf{b} \in \hat{\mathcal{M}}$ and
 260 $\mathbf{x}_b \in N_{\varepsilon'}(\mathbf{x}_0)$.

261 We can further assume $\hat{\mathcal{M}}$ to be sufficiently small so that, for all $\mathbf{b} \in \hat{\mathcal{M}} \subset \tilde{\mathcal{M}}$,

$$\|(\mathbf{f}_x(\mathbf{y})^\dagger - \mathbf{f}_x(\mathbf{z})^\dagger)(\mathbf{f}(\mathbf{z}) - \mathbf{b})\| \leq \frac{1}{2}\|\mathbf{y} - \mathbf{z}\|$$

262 whenever $\mathbf{y} \in N_\varepsilon(\mathbf{x}_0)$ and $\mathbf{z} \in N_{\varepsilon'}(\mathbf{x}_0)$, making

$$\|\mathbf{y}_0 - \mathbf{x}_b\| < \frac{1}{2\zeta\gamma} = \frac{1 - \frac{1}{2}}{\zeta\gamma}$$

263 for every $\mathbf{y}_0 \in N_{\varepsilon'}(\mathbf{x}_0)$. Furthermore, the inequality $\|\mathbf{x} - \mathbf{x}_b\| < \|\mathbf{y}_0 - \mathbf{x}_b\|$ with
 264 $\mathbf{y}_0 \in N_{\varepsilon'}(\mathbf{x}_0)$ implies

$$\begin{aligned} \|\mathbf{x} - \mathbf{x}_0\| &\leq \|\mathbf{x} - \mathbf{x}_b\| + \|\mathbf{x}_b - \mathbf{x}_0\| \\ &< \|\mathbf{y}_0 - \mathbf{x}_b\| + \|\mathbf{x}_b - \mathbf{x}_0\| \\ &\leq \|\mathbf{y}_0 - \mathbf{x}_0\| + \|\mathbf{x}_0 - \mathbf{x}_b\| + \varepsilon' \\ &< \varepsilon' + \varepsilon' + \varepsilon' = \varepsilon. \end{aligned}$$

265 Namely $\{\mathbf{x} \in \mathcal{V} \mid \|\mathbf{x} - \mathbf{x}_b\| < \|\mathbf{y}_0 - \mathbf{x}_b\|\} \subset N_\varepsilon(\mathbf{x}_0)$. As a result, the Gauss-Newton
 266 iteration

$$\mathbf{y}_{k+1} = \mathbf{y}_k - \mathbf{f}_x(\mathbf{y}_k)^\dagger (\mathbf{f}(\mathbf{y}_k) - \mathbf{b}) \quad \text{for } k = 0, 1, \dots \quad (14)$$

267 converges to \mathbf{x}_b from every $\mathbf{y}_0 \in N_{\varepsilon'}(\mathbf{x}_0)$. This least squares solution \mathbf{x}_b is
 268 unique for every $\mathbf{b} \in \hat{\mathcal{M}}$. In fact, another least squares solution $\tilde{\mathbf{x}}_b$ to the
 269 equation $\mathbf{f}(\mathbf{x}) = \mathbf{b}$ implies the Gauss-Newton iteration (14) converges to \mathbf{x}_b from
 270 the initial iterate $\mathbf{y}_0 = \tilde{\mathbf{x}}_b$. On the other hand, the Gauss-Newton iteration starting
 271 from the critical point $\tilde{\mathbf{x}}_b$ stays at $\tilde{\mathbf{x}}_b$, leading to $\tilde{\mathbf{x}}_b = \mathbf{x}_b$.

272 We have now established a mapping $\mathbf{b} \mapsto \mathbf{x}_b$ from the domain $\hat{\mathcal{M}}$ to the manifold
 273 Π . We claim this mapping is Lipschitz continuous. Let $\mathbf{x}_d \in N_{\varepsilon'}(\mathbf{x}_0)$ be the least
 274 squares solution to the equation $\mathbf{f}(\mathbf{x}) = \mathbf{d}$ and perform one step of the Gauss-Newton
 275 iteration

$$\mathbf{x}_1 = \mathbf{x}_b - \mathbf{f}_x(\mathbf{x}_b)^\dagger (\mathbf{f}(\mathbf{x}_b) - \mathbf{d})$$

276 for the equation $\mathbf{f}(\mathbf{x}) = \mathbf{d}$ from the initial iterate \mathbf{x}_b . Using the identity
 277 $\mathbf{x}_b = \mathbf{x}_b - \mathbf{f}_x(\mathbf{x}_b)^\dagger (\mathbf{f}(\mathbf{x}_b) - \mathbf{b})$, we have

$$\|\mathbf{x}_1 - \mathbf{x}_b\| \leq \|\mathbf{f}_x(\mathbf{x}_b)^\dagger\| \|\mathbf{b} - \mathbf{d}\|.$$

278 On the other and, Lemma 5.1 implies that there is a $\mu \leq \frac{1}{2} + \zeta\gamma\|\mathbf{x}_b - \mathbf{x}_d\|$ such
 279 that

$$\|\mathbf{x}_1 - \mathbf{x}_d\| \leq \mu \|\mathbf{x}_b - \mathbf{x}_d\|.$$

280 Consequently

$$\begin{aligned} \|\mathbf{x}_d - \mathbf{x}_b\| &\leq \|\mathbf{x}_b - \mathbf{x}_1\| + \|\mathbf{x}_1 - \mathbf{x}_d\| \\ &\leq \|\mathbf{f}_x(\mathbf{x}_b)^\dagger\| \|\mathbf{b} - \mathbf{d}\| + \mu \|\mathbf{x}_b - \mathbf{x}_d\| \end{aligned}$$

281 and thus

$$\|\mathbf{x}_b - \mathbf{x}_d\| \leq \frac{1}{1 - \mu} \|\mathbf{f}_x(\mathbf{x}_b)^\dagger\| \|\mathbf{b} - \mathbf{d}\|. \quad (15)$$

282 The asymptotic upper bound of μ is $\frac{1}{2}$ when \mathbf{d} approaches to \mathbf{b} and μ can be
 283 arbitrarily small if \mathbf{b} is close to the manifold Π .

284 Let Ω be the tubular neighborhood as the union of all those $\hat{\mathcal{M}}$ over all \mathbf{z}_0 . The
 285 theorem is thus proved. \square

7 Geometry of hypersensitive problems

Attributed to Hadamard, the notion of *well-posed problem* requires a mathematical model to have a solution that satisfies existence, uniqueness and some kind of continuity with respect to data. For the problem to be feasible in numerical computation, its solution needs to be Lipschitz continuous so that the sensitivity is bounded with respect to data perturbation and round-off. The version of Tubular Neighborhood Theorem in Theorem 4.1 is specifically formulated and established to ensure that the (numerical) least squares problem of computing the nearest point on a complex analytic manifold from any point in the tubular neighborhood is a well-posed problem.

A problem whose solution lacks one of the existence, uniqueness or continuity is called an *ill-posed problem*. In particular, if the solution is not Lipschitz continuous with respect to data, we say the problem is *hypersensitive* with a condition number infinity. A typical textbook example in numerical analysis is the root-finding problem of polynomials in the form of solving the equation

$$a_0 + a_1 z + a_2 z^2 + \cdots + a_n z^n = 0 \quad (16)$$

with $a_n \neq 0$. If the roots of the polynomial $p(z) = a_0 + a_1 z + a_2 z^2 + \cdots + a_n z^n$ are all simple, every solution z_* of (16) has a finite condition number $\frac{1}{|p'(z_*)|}$ and is thus Lipschitz continuous with respect the coefficient vector (a_0, a_1, \dots, a_n) serving as data. When multiple roots are present, however, the Lipschitz continuity is lost and the root-finding problem (16) becomes hypersensitive.

Geometrically, every such a polynomial p is a point in the vector space \mathbb{P}_n that consists of polynomials with degrees up to n . If p has k roots of multiplicities ℓ_1, \dots, ℓ_k respectively where $\ell_1 + \cdots + \ell_k = n$, then p is on the complex analytic manifolds $\mathcal{F}_{\ell_1 \dots \ell_k}$ defined in (4) with codimension $n - k$. The root-finding problem for $p(z) = 0$ is hypersensitive except for p to be on the trivial factorization manifold $\mathcal{F}_{1 \dots 1}$, i.e. $k = n$, $\ell_1 = \cdots = \ell_n = 1$ and $\text{codim}(\mathcal{F}_{\ell_1 \dots \ell_n}) = n - n = 0$. We can quantify the hypersensitivity of a polynomial p in terms of root-finding by defining the *singularity* of $p \in \mathcal{F}_{\ell_1 \dots \ell_n}$ as $\text{codim}(\mathcal{F}_{\ell_1 \dots \ell_n})$. We also say the *singularity* of the complex analytic manifold $\mathcal{F}_{\ell_1 \dots \ell_n}$ is $n - k$. Notice that the trivial manifold $\mathcal{F}_{1 \dots 1}$ of singularity zero is open and dense in \mathbb{P}_n , leading to the observation:

Nonsingular problems are open and dense in the problem space. A problem is hypersensitive only if the data point resides on a manifold of nonzero singularity (codimension).

When a simple root z_* of p is ill-conditioned, namely $p'(z_*) \approx 0$. The polynomial p is clearly near a polynomial with multiple roots and is thus near a manifold $\mathcal{F}_{\ell_1 \dots \ell_n}$ of nonzero singularity $n - k$, implying:

A problem is ill-conditioned only if the data point is near a manifold of nonzero singularity (codimension).

323 Consider the holomorphic mapping

$$\begin{aligned} \mathbf{f} & : & \mathbb{C}^{k+1} & \longrightarrow & \mathbb{P}_n \\ & & (z_0, z_1, \dots, z_k) & \longmapsto & z_0(x - z_1)^{\ell_1} \dots (x - z_k)^{\ell_k}. \end{aligned}$$

324 For every $p \in \mathcal{F}_{\ell_1 \dots \ell_k}$, there is a $\mathbf{z} \in \mathbb{C}^{k+1}$ such that $\mathbf{f}(\mathbf{z}) = p$ and there are
 325 neighborhoods Λ and Δ of p and \mathbf{z} , respectively, along with a holomorphic
 326 mapping $\mathbf{g} : \Lambda \rightarrow \Delta$ that is the inverse of \mathbf{f} in Δ . As a result, the mapping
 327 \mathbf{f} serves as the local diffeomorphism for the complex analytic manifold $\mathcal{F}_{\ell_1 \dots \ell_k}$. By
 328 the proof of the Tubular Neighborhood Theorem, we have

$$\|\mathbf{z} - \tilde{\mathbf{z}}\|_2 \leq \frac{1}{1 - \mu} \|\mathbf{f}_{\mathbf{z}}(\mathbf{z})^\dagger\| \|p - \tilde{p}\|$$

329 for $\tilde{p} = \mathbf{f}(\tilde{\mathbf{z}})$ near p and $\mu \rightarrow 0$ when \tilde{p} approaches p on the manifold
 330 $\mathcal{F}_{\ell_1 \dots \ell_k}$. Namely, even though the roots of the polynomial p are hypersensitive to
 331 *arbitrary perturbations* on p , they are in the solution of the equation $\mathbf{f}(\mathbf{z}) = p$
 332 that is Lipschitz continuous for p constrained on the manifold $\mathcal{F}_{\ell_1 \dots \ell_k}$, validating
 333 Kahan's observation in 1972 [9]:

334 *A singular problem is not hypersensitive if the perturbation is con-*
 335 *strained such that the data point stays on its native manifold.*

336 Perhaps the most crucial observation of hypersensitive problems is the geometric
 337 explanation of the hypersensitivity: A problem is hypersensitive because it resides
 338 on a complex analytic manifold of nonzero singularity (codimension). Due to the
 339 dimension deficit, an arbitrary perturbation generically pushes the data (point) away
 340 from the native manifold and thus changes the solution structure completely no matter
 341 how tiny the perturbation is. Looking deeper, however, we need to point out that the
 342 solution changes from tiny perturbations in one and only one direction in reducing
 343 the singularity. Consider a polynomial, say

$$p(x) = (x - 1)^2(x - 2)^3 \in \mathcal{F}_{23}$$

344 of singularity $5 - 2 = 3$. It is arbitrarily close to the infinitely many polynomials of
 345 different multiplicity structures such as

$$p_\varepsilon(x) = (x - 1 - \varepsilon)(x - 1 + \varepsilon)(x - 2)^3 \in \mathcal{F}_{113}$$

346 of singularity $5 - 3 = 2$ when $\varepsilon \rightarrow 0$. Clearly the manifold $\mathcal{F}_{23} \subset \overline{\mathcal{F}_{113}}$. We say
 347 a manifold Π is *embedded* in the manifold Σ , denoted by $\Pi \hookrightarrow \Sigma$ if the former is in
 348 the closure of the latter. It is common for hypersensitive problems to form complex
 349 analytic manifolds that entangle and form *strata* in which a singular manifold is
 350 embedded in manifolds of lower singularities. Figure 3 shows the stratification of
 351 manifolds of degree-5 polynomials.

352 As illustrated in Figure 3, the manifold \mathcal{F}_{23} is embedded in \mathcal{F}_{113} , \mathcal{F}_{122} , \mathcal{F}_{1112} and
 353 the nonsingular \mathcal{F}_{11111} . As a result, a small perturbation Δp on $p \in \mathcal{F}_{23}$ can push
 354 the polynomial p to $p + \Delta p$ into those manifolds of lower singularities. There is a
 355 minimum distance δ from p to \mathcal{F}_{14} and \mathcal{F}_5 of equal or higher singularities and
 356 the perturbation Δp must have a magnitude at least δ to change the multiplicity
 357 structure of p to equal or higher singularities. Thus:

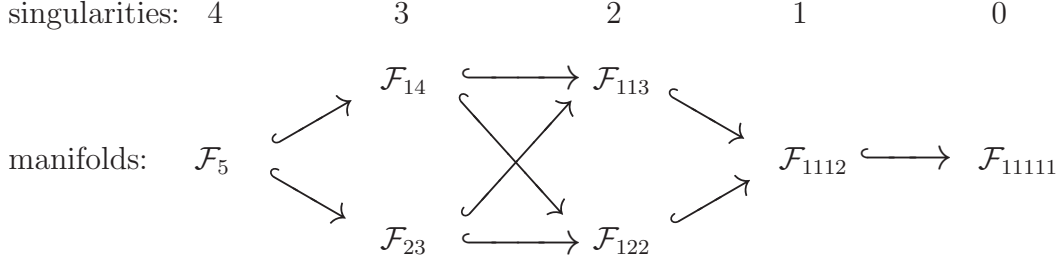


Figure 3: Stratification of manifolds of degree-5 polynomials and their singularities.

358 *A small perturbation on the data of a hypersensitive problem can only*
 359 *reduce the singularity and never increase it.*

360 This unidirectional perturbation property of singularity is the most crucial property
 361 for numerical computation and provides a mechanism for identifying the manifold on
 362 which the problem resides. We define the distance of a point \mathbf{z} from a manifold Π
 363 in a normed vector space as

$$\text{dist}(\mathbf{z}, \Pi) := \inf_{\mathbf{y} \in \Pi} \|\mathbf{y} - \mathbf{z}\|.$$

364 Such a distance is zero if and only if $\mathbf{z} \in \bar{\Pi}$. For instance, a polynomial $p \in \mathcal{F}_{23}$
 365 means that p has a zero distance to manifolds \mathcal{F}_{23} , \mathcal{F}_{113} , \mathcal{F}_{122} , \mathcal{F}_{1112} and \mathcal{F}_{11111} .
 366 Among those manifolds, the native manifold \mathcal{F}_{23} distinguishes itself as the one with
 367 the highest singularity

$$\text{codim}(\mathcal{F}_{23}) = \max_{\substack{\text{dist}(p, \mathcal{F}_{\ell_1 \dots \ell_k}) = 0, \\ \ell_1 + \dots + \ell_k = 5}} \text{codim}(\mathcal{F}_{\ell_1 \dots \ell_k})$$

368 leading to the observation:

369 *The native manifold of a hypersensitive problem is identifiable as the*
 370 *manifold with the highest singularity among all manifolds with a zero dis-*
 371 *tance to the problem.*

372 In practical computation, a hypersensitive problem is likely represented by imperfect
 373 empirical data while the precise values are unknown. Even if the exact data are
 374 known, numerical algorithms inevitably introduce round-off errors during the solution
 375 process. Continuing the example of the root-finding problem, even if the polynomial
 376 $p \in \mathcal{F}_{23}$ is represented by the empirical data point $p + \Delta p$ where Δp is arbitrary,
 377 the native manifold \mathcal{F}_{23} is still identifiable as long as the magnitude $\|\Delta p\|$ of the
 378 perturbation is not too large. Let

$$\begin{aligned} \delta &= \min_{\substack{\text{dist}(p, \mathcal{F}_{\ell_1 \dots \ell_k}) > 0, \\ \ell_1 + \dots + \ell_k = 5}} \text{dist}(p, \mathcal{F}_{\ell_1 \dots \ell_k}) \\ &= \min \{ \text{dist}(p, \mathcal{F}_{14}), \text{dist}(p, \mathcal{F}_5) \} \end{aligned}$$

379 that is nonzero. Assume $\|\Delta p\| < \frac{\delta}{2}$. The empirical data point $\tilde{p} = p + \Delta p$ can
 380 land on any one of the manifolds \mathcal{F}_{23} , \mathcal{F}_{113} , \mathcal{F}_{122} , \mathcal{F}_{1112} and \mathcal{F}_{11111} but not \mathcal{F}_{14}
 381 or \mathcal{F}_5 . The native manifold \mathcal{F}_{23} of the underlying problem p is still the manifold
 382 of the highest singularity among all the manifolds within a distance $\frac{\delta}{2}$ of the data.
 383 Namely

$$\text{codim}(\mathcal{F}_{23}) = \max_{\substack{\text{dist}(p+\Delta p, \mathcal{F}_{\ell_1 \dots \ell_k}) < \frac{\delta}{2}, \\ \ell_1 + \dots + \ell_k = 5}} \text{codim}(\mathcal{F}_{\ell_1 \dots \ell_k})$$

384 The problem of identifying the native manifold \mathcal{F}_{23} by maximizing the singularity
 385 within $\frac{\delta}{2}$ of the data point $p + \Delta p$ is clearly a well posed problem.

386 *The native manifold of the underlying hypersensitive problem is iden-*
 387 *tifiable from its empirical data as the manifold of the highest singularity*
 388 *among all manifolds with a distance below a proper error tolerance from*
 389 *the data point.*

390 Assuming the data point $\tilde{p} = p + \Delta p$ is reasonably close to p and inside the tubular
 391 neighborhood of the native manifold \mathcal{F}_{23} , the Tubular Neighborhood Theorem
 392 ensures that there is a unique projection from \tilde{p} to a point $\hat{p} = \mathbf{f}(\hat{\mathbf{z}})$ on the
 393 manifold \mathcal{F}_{23} and $\hat{\mathbf{z}}$ is an accurate numerical solution to the underlying problem
 394 and the solution error is in the same order of the data error.

395 The seven geometric observations can be rigorously proved for polynomial root-finding
 396 problem, polynomial greatest common divisor problem in both univariate and mul-
 397 tivariate cases, multivariate polynomial factorization problem, etc. They are likely
 398 true for many other hypersensitive problems such as the problem of computing the
 399 Jordan Canonical Form of matrices.

400 We can reiterate those observations on the rank-revealing problem. Let A be a
 401 matrix of size, say, 8×5 . The rank-revealing problem of A consists computing the
 402 rank and subspaces such as ranges and kernels. The matrix A is on the rank manifold
 403 $\mathcal{C}_r^{8 \times 5}$ where r is the rank of A . The singularity $\text{codim}(\mathcal{C}_r^{8 \times 5}) = (8 - r)(5 - r)$.
 404 The rank-revealing problem of A is hypersensitive if and only if the singularity is
 405 nonzero (i.e. $\text{rank}(A) < 5$). When A is of full-rank, the rank-revealing problem
 406 is ill-conditioned only if A is near rank-deficient matrices so that the condition
 407 number of A is large. When A is rank-deficient with $r < 5$, tiny arbitrary
 408 perturbations on A alter the rank and subspaces substantially. If the perturbation
 409 ΔA is constrained so that $A + \Delta A$ stays on the manifold $\mathcal{C}_r^{8 \times 5}$, however, the rank
 410 and subspaces are not hypersensitive at all by Wedin's Theorem [12, Theorem 4.1, p.
 411 260]. It is well-known in linear algebra that a tiny perturbation can only increase the
 412 rank and never reduce it, or equivalently can only reduces the singularity and never
 413 increase it. The rank manifolds in the vector space $\mathbb{C}^{8 \times 5}$ form a simple strata

$$\begin{array}{ccccccccc} \text{singularities:} & 40 & & 28 & & 18 & & 10 & & 4 & & 0 \\ \text{manifolds} & \mathcal{C}_0^{8 \times 5} & \hookrightarrow & \mathcal{C}_1^{8 \times 5} & \hookrightarrow & \mathcal{C}_2^{8 \times 5} & \hookrightarrow & \mathcal{C}_3^{8 \times 5} & \hookrightarrow & \mathcal{C}_4^{8 \times 5} & \hookrightarrow & \mathcal{C}_5^{8 \times 5}. \end{array}$$

414 If the rank of A is, say, 3, then A is of zero distance to manifolds $\mathcal{C}_3^{8 \times 5}$, $\mathcal{C}_4^{8 \times 5}$
 415 and $\mathcal{C}_5^{8 \times 5}$ and the native manifold $\mathcal{C}_3^{8 \times 5}$ distinguishes itself as having the highest

416 singularity among them. Even if A is only known thorough its empirical data as
 417 $\tilde{A} = A + \Delta A$, the native manifold $\mathcal{C}_3^{8 \times 5}$ still distinguishes itself as having the highest
 418 singularity among nearby rank-manifolds.

419 8 Regularization of hypersensitive problems

420 When we face the root-finding problem of a polynomial, say $\tilde{p}(x)$ in (1), what we
 421 really want are the roots of the hidden underlying polynomial $p(x)$ in (2). with a
 422 data error bound about 2×10^{-3} . However, the polynomial p is hypersensitive for
 423 the root-finding problem in conventional sense. As shown in Figure 1, it makes little
 424 sense to solve the equation $\tilde{p}(x) = 0$ since the solutions bear almost no resemblance
 425 to the solutions we really expect.

426 The objective in practical computation is to *find the solution of the underlying problem*
 427 *accurately from empirical data* with an accuracy the data deserve. In our simulated
 428 example, the goal is to find roots of p in (2) accurately using the data polynomial
 429 such as \tilde{p} in (1) with an accuracy in the same order of the data accuracy without
 430 knowing p beyond the error bound $\|p - \tilde{p}\| \leq 2 \times 10^{-3}$. For this purpose, we
 431 have no choice but to abandon the meaningless problem of “solving $\tilde{p}(x) = 0$ ” and
 432 modify it to a numerically feasible problem. This problem modification process for
 433 an ill-posed problem is call *regularization*.

434 Ideally, the regularized problem should have the identical solution at the specific un-
 435 perturbed data of the original problem and the solution becomes Lipschitz continuous
 436 with respect to data in a neighborhood of the intended solution.

437 From the geometry, the underlying polynomial p is on the manifold $\mathcal{F}_{1,4}$ that is
 438 identifiable as having the highest singularity among manifolds nearby, and roots are
 439 not hypersensitive with respect to the polynomials on that manifold. We can natu-
 440 rally regularize the root-finding problem as calculating a specific type of *regularized*
 441 *roots* that are exact roots of the nearest polynomial on the factorization manifold of
 442 highest singularity passing through a neighborhood of the data polynomial.

443 **Definition 8.1 (Regularized Roots)** *Let \tilde{p} be a polynomial of degree n along*
 444 *with an error tolerance ε . The regularized roots of \tilde{p} within ε are exact roots of*
 445 *a polynomial \hat{p} that is on the factorization manifold $\mathcal{F}_{\ell_1 \dots \ell_k}$ where*

$$\begin{aligned} \text{codim}(\mathcal{F}_{\ell_1 \dots \ell_k}) &= \max_{\substack{\text{dist}(\tilde{p}, \mathcal{F}_{\lambda_1 \dots \lambda_l}) < \varepsilon, \\ \lambda_1 + \dots + \lambda_l = n}} \text{codim}(\mathcal{F}_{\lambda_1 \dots \lambda_l}) \\ \|\tilde{p} - \hat{p}\| &= \min_{u \in \mathcal{F}_{\ell_1 \dots \ell_k}} \|\tilde{p} - u\| \end{aligned}$$

446 **Theorem 8.2 (Fundamental Theorem of Regularized Roots)** *Let p be a*

447 polynomial of degree n with roots z_1, \dots, z_k of multiplicities ℓ_1, \dots, ℓ_k respec-
 448 tively. The following properties of regularized roots hold.

449 (i) The notion of regularized roots generalizes exact roots: There is a $\mu > 0$
 450 such that the regularized roots of p within any $\varepsilon \in (0, \mu)$ are identical to the
 451 exact roots of p with same respective multiplicities.

452 (ii) The regularized root problem is well-posed: There is a neighborhood Ω_p of
 453 p in which every polynomial \tilde{p} serving as empirical data of p is associated
 454 with an interval $(\delta_{\tilde{p}}, \theta_{\tilde{p}})$ with $0 \leq \delta_{\tilde{p}} \leq \|p - \tilde{p}\|$ such that the regularized roots
 455 of \tilde{p} within $\varepsilon \in (\delta_{\tilde{p}}, \theta_{\tilde{p}})$ uniquely exist as $\tilde{z}_1, \dots, \tilde{z}_k$ of the same multiplicities
 456 ℓ_1, \dots, ℓ_k respectively and are Lipschitz continuous with respect to \tilde{p} in the
 457 sense that, for every $\tilde{q} \in \Omega_p$, the regularized roots $\tilde{y}_1, \dots, \tilde{y}_k$ of \tilde{q} with
 458 multiplicities ℓ_1, \dots, ℓ_k respectively satisfy

$$\|(\tilde{y}_1, \dots, \tilde{y}_k) - (\tilde{z}_1, \dots, \tilde{z}_k)\| \leq \eta_p \|\tilde{p} - \tilde{q}\|.$$

459 where η_p is a constant depends on p and ε .

460 (iii) Regularized roots are backward accurate: For every $\tilde{p} \in \Omega_p$, the regularized
 461 roots of \tilde{p} within $\varepsilon \in (\delta_{\tilde{p}}, \theta_{\tilde{p}})$ are exact roots of a polynomial \hat{p} with

$$\|\tilde{p} - \hat{p}\| \leq \|\tilde{p} - p\| < \varepsilon.$$

462 (iv) The regularized roots of data polynomial \tilde{p} approximate the exact roots of the
 463 underlying polynomial p with an accuracy $O(\|p - \tilde{p}\|)$.

464 Quite obviously, the well-posedness assertion (ii) in the theorem is the corollary of
 465 the Tubular Neighborhood Theorem as established in Theorem 4.1.

466 This theorem is a special case of the Numerical Factorization Theorem in [14] with
 467 a rigorous and detailed proof referencing the Tubular Neighborhood Theorem in an
 468 unpublished technical report by this author. Theorem 4.1 completes the work of
 469 numerical factorization in [14] and finishes the proof of Fundamental Theorem of
 470 Regularized Roots.

471 This regularization follows “three-strikes principles” that have been applied effectively
 472 for striking out the hypersensitivity of singular algebraic problems:

473 *Backward stability:* The regularized numerical solution at given data is the exact
 474 solution at a specific data point within an error tolerance.

475 *Maximum singularity:* The specific data point is on a particular manifold of the
 476 highest singularity among all manifolds within the error tolerance of the given
 477 data.

478 *Minimum distance:* The specific data point is the nearest point on the particular
 479 manifold to the given data.

480 A hypersensitive problem such as the roots of the polynomial p in (2) is on a complex
 481 analytic manifold such as \mathcal{F}_{41} . Solving such problems from empirical data is possible
 482 only if the data are reasonably sound so that the data point is still in the tubular
 483 neighborhood.

9 Numerical solution of the regularized problem

In a nutshell, the geometric regularization can be summarized in one sentence:

The regularized solution of a hypersensitive problem is the exact solution of a nearby problem that is the point on the manifold of maximum singularity with the minimum distance from the data.

A natural approach for finding the regularized solution is a two staged process:

Stage I. Within the error tolerance of the data, find the nearby complex analytic manifold of the highest singularity.

Stage II. Solve the least squares problem that minimizes the distance from the data to the manifold.

We continue using the polynomial root-finding problem for elaboration of the algorithmic engineering. Given the data polynomial \tilde{p} in (1), the data error is at the third digits of the coefficients so the error tolerance can be set at 10^{-3} (more to that later).

The underlying polynomial p in (2) is on the manifold \mathcal{F}_{41} if and only if the polynomial pair (p, p') belongs to the complex analytic manifold $\mathcal{P}_{5,4}^3$ since

$$\gcd(p, p') = (x - 1)^3.$$

The singularity of $\mathcal{P}_{5,4}^3$ is identical to the degree 3 of the greatest common divisor. Write $p(x) = (x - 1)^3 \cdot ((x + 1)(x + \frac{2}{3}))$ and

$$3 = 5 - \deg((x - 1)(x + \frac{2}{3})) = n - k$$

where $n = \deg(p)$ and k is the number of roots. It is easy to see that finding the maximal singularity factorization polynomial near \tilde{p} is equivalent to finding the maximal singularity GCD manifold near the polynomial pair (\tilde{p}, \tilde{p}') .

The polynomial GCD problem is also hypersensitive and can be regularized following the same three-strikes principles [18]. Using the numerical GCD finder `PolynomialGCD` in the package `NACLAB`², we can calculate the *regularized* greatest common divisor within the error tolerance 10^{-3} by the following calls:

```
>> [u,v,w,res,gcond]=PolynomialGCD(p,q,1e-3)
u =
0.89678776072183 + 2.68248022385017*x + 2.6777673064563*x^2 + 0.891138887558938*x^3
v =
0.743447961578301 + 1.86510244285362*x + 1.12216238113407*x^2
```

²<http://homepages.neiu.edu/~nac1ab>

```

515
516 w =
517 4.08901144885262 + 5.61078724773379*x
518
519 res =
520 3.122507798584681e-05
521
522 gcond =
523 5.630826706824984e+02

```

524 The results can be interpreted as follows: The polynomial pair (\tilde{p}, \tilde{p}') has a regu-
525 larized greatest common divisor \hat{u} and co-factors \hat{v} and \hat{w} such that

$$\tilde{p} \approx \hat{p} = \hat{u} \hat{v}, \quad \tilde{q} \approx \hat{q} = \hat{u} \hat{w}.$$

526 The regularized co-factor $\hat{v} \approx 0.7435 + 1.8651x - 1.1222x^2$ of \tilde{p} is of degree 2
527 with roots approximately 0.9987 and 0.6634. Furthermore, rounding up $\frac{\hat{w}(x)}{\hat{v}'(x)}$ to
528 the nearest integer at x values -0.9987 and -0.6634 yields 4 and 1 respectively.
529 As a result, the most singular manifold near \tilde{p} within the error tolerance 10^{-3} is
530 identified as \mathcal{F}_{41} with roots approximately -0.9987 and -0.6634 of multiplicity
531 4 and 1 respectively, completing Stage I of the process.

532 Every polynomial in \mathcal{F}_{41} is an image of the holomorphic mapping

$$\mathbf{f} : \mathbb{C}^3 \longrightarrow \mathbb{P}_5$$

$$(a, z_1, z_2) \longmapsto a(x - z_1)^4(x - z_2)$$

533 Thus Stage II of the computation can be carried out by solving for the least squares
534 solution of the system

$$\mathbf{f}(a, z_1, z_2) = \tilde{p}$$

535 by the Gauss-Newton iteration with the initial iterate

$$(a_0, z_1^{(0)}, z_2^{(0)}) = (1, -0.9987, -0.6634).$$

536 The complete algorithm for computing the regularized roots is implemented in our
537 Matlab package NACLAB as the module PolynomialFactor. The numerical result
538 can be obtained with a simple call:

```

539 >> [P,res,fcond] = PolynomialFactor(p,1e-3,'row')
540
541 P =
542 (1.00006771643903) * (x+0.999940850183)^4 * (x+0.666823116178)
543
544 res =
545 3.624323682583038e-04
546
547 fcond =
548 7.938647029187908

```

549 Namely, the numerical solution of the regularized factorization is

$$1.00006771643903(x + 0.999940850183)^4(x + 0.666823116178)$$

550 with a residual 3.62×10^{-4} and a regularized condition number 7.94 measuring the
551 sensitivity.

552 The computed regularized roots are -0.99994 and -0.66682, approximating the exact
553 roots -1 and $-\frac{2}{3}$ with the precise multiplicities 4 and 1 respectively. Considering
554 the data error bound 10^{-3} , this is the accuracy the data deserve. The module
555 `PolynomialFactor` also reports the residual as 3.62×10^{-4} that is the distance from
556 the data \tilde{p} to the manifold \mathcal{F}_{41} serving as the backward accuracy of the computing
557 result.

558 Most importantly, the sensitivity measure 7.94 indicates that computing the reg-
559 ularized roots is highly stable even though the conventional root-finding problem is
560 highly sensitive. As a result, we achieved a complete geometrical regularization of
561 the ill-posed problem by applying the Tubular Neighborhood Theorem.

562 **Acknowledgment.** The author is indebted to his former colleague Marian Gidea
563 for a casual chat in the department copy room several years ago. Knowing nothing
564 about differential geometry beyond roughly what a manifold was, I asked: If a point
565 is near a smooth manifold, isn't it intuitive that the projection of the point to the
566 manifold is unique? His response: It is just the Tubular Neighborhood Theorem.
567 The theorem is, of course, in his book [2]. That chat precipitated several papers
568 including this one.

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